

## System Software Research (1)



## Towards Next-Generation Supercomputing

## Performance Analysis and Runtimes

## **Performance Metrics and Tools**

We've installed following performance monitoring tools into TSUBAME2.5 and can get performance data of real applications in following metrics.

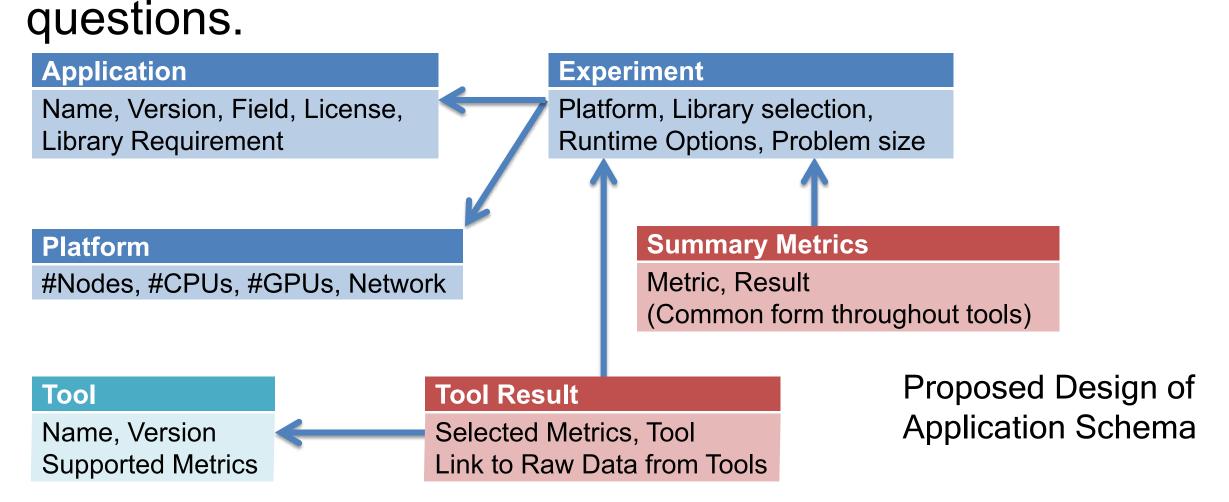
Tool	Type	Metrics
Scalasca	Profiler	Time, Visit, MPI Communication, PAPI
Vampir	Tracer	Time, MPI Comm., GPU Comm., PAPI
Exana	Tracer	Memory Access (Instruction Level)
PAPI	Library	Instruction Mix, Main Memory Access

We're giving seminar for those performance tools for TSUBAME users.

## Performance Repository Schema

We're building application performance repository to know following things through performance analysis.

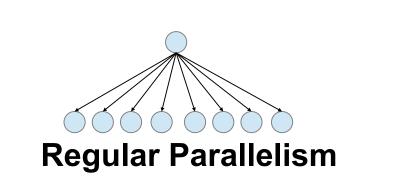
- Which machine can solve my application well?
- Which application are suitable for our supercomputer? We collect various performance data from many supercomputers and correlate them to answer these

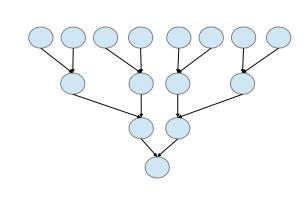


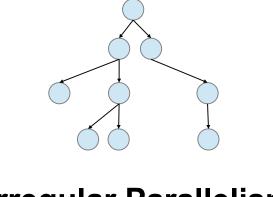
This is collaborative work with Future Technologies Group, Oak Ridge National Laboratory.

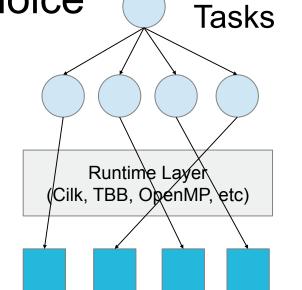
## **Task Parallel Runtimes**

Support for many types of control makes them a popular choice









**Divide and Conquer** 

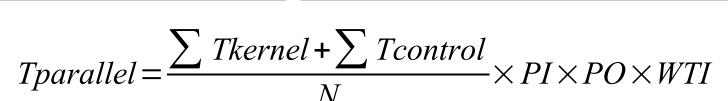
Irregular Parallelism

#### Common issues in runtime scalability:

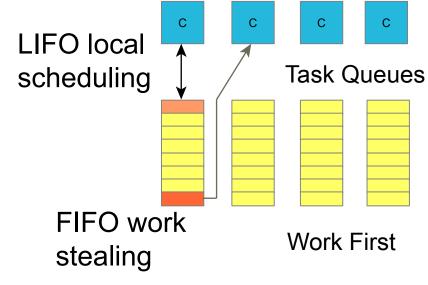
1. Runtime System overheads Causes parallelization overhead (PO) 2. Scheduling Constraints Causes parallel idleness (PI)

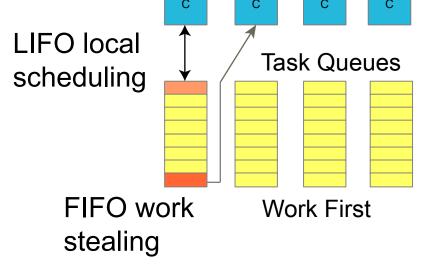
3. Resource contention Causes Work Time Inflation (WTI)

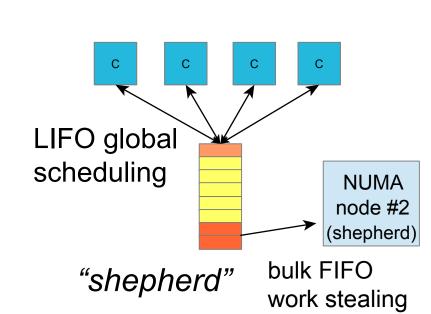
 $Tserial = \sum Tkernel + \sum Tcontrol$ 



## Case Study: scalability of recursive MatMul





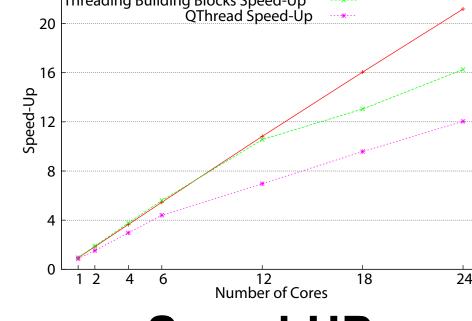


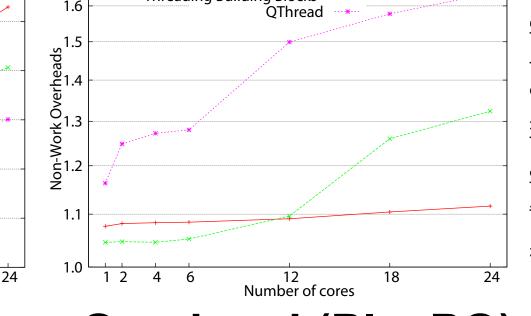
**MassiveThreads** 

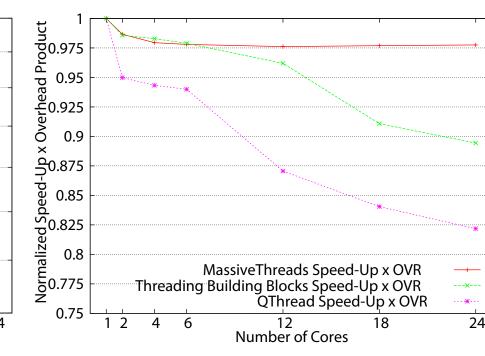
**Thread Building Blocks** 

**QThread** 

**Experiment Details:** 4-socket Intel Xeon E7-4807 (1.86GHz). Matrix size 4096<sup>2</sup> elem (64MB) ding Building Blocks Speed-Up QThread Speed-Up







Speed-UP

Overhead (PI x PO)

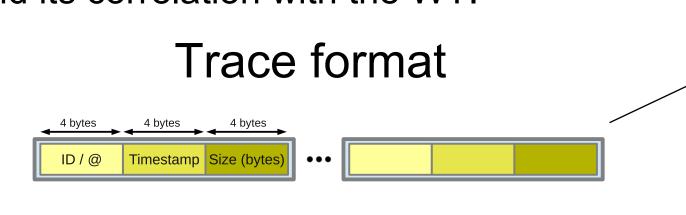
Speed-Up x OVR **Product (1/WTI)** 

## Data Reuse and Heterogeneity

## **KRD** metric

### **Kernel Reuse Distance**

- Reuse distance of Kernel data accesses
- Coarse Grain to reduce overheads
- Tool for the analysis of LLC performance and its correlation with the WTI



## Data Access Trace Synchronized/Merged Trace generation MAIN MEMORY

## Heterogeneous Scheduling

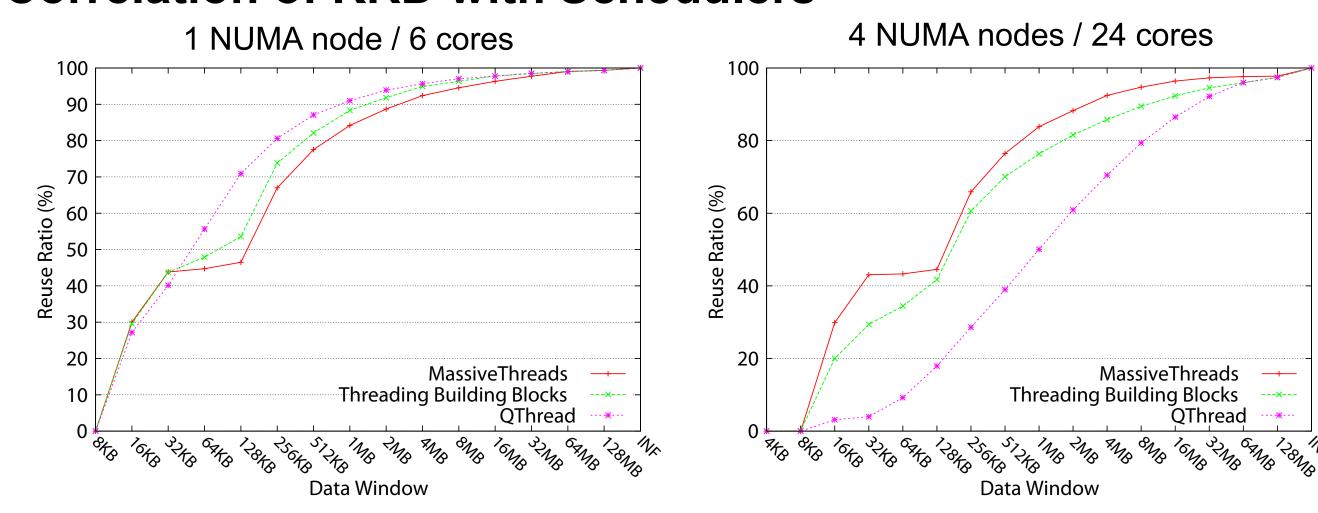
# Tree construction

## **Case Study: Fast Multipole Method**

- 1) Complex inter-task dependencies
- 2) Heterogeneous task runtimes
- 3) Input dependent control flow
- 4) Mix of compute/memory bound tasks

## Where and when to schedule each task on a CPU/GPU system?

#### **Correlation of KRD with Schedulers**



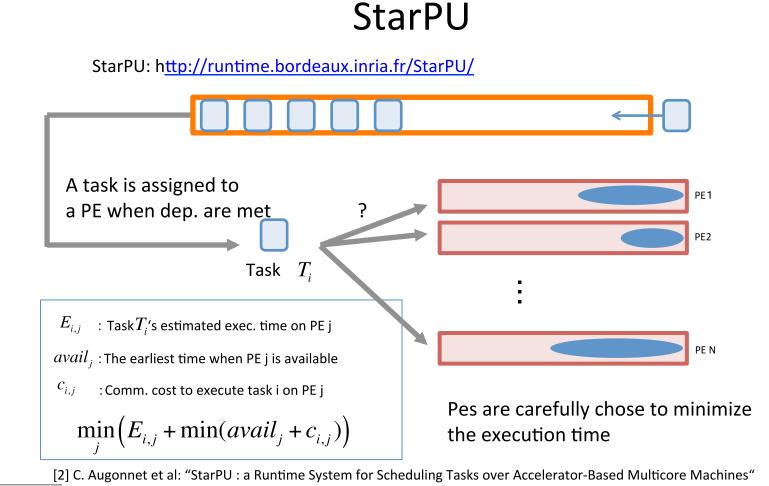
#### Approach: Use dynamic task scheduler

Update particle informatio

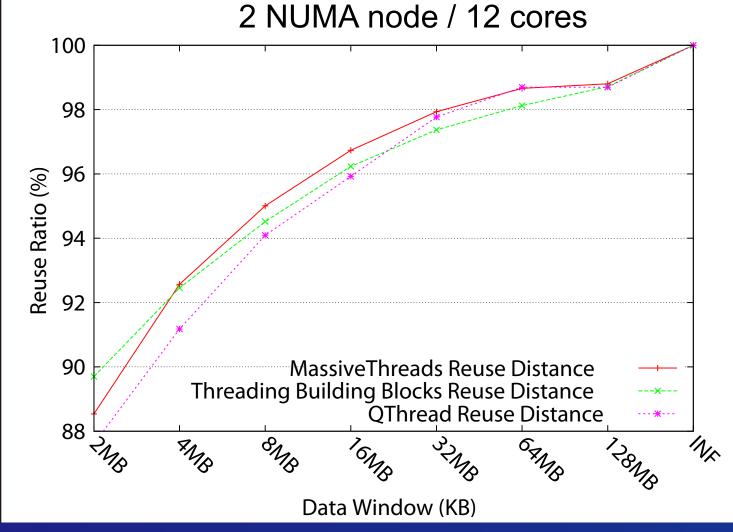
Comparison of StarPU with:

1) All-GPU: All tasks on GPU 2) Simple Hybrid, P2P on GPU and remaining phases on CPU 3) Optimal scheduling, based on design space exploration

A dynamic task scheduling engine:



#### **Correlation of KRD with Hardware Metrics**





Global LIFO increases locality in single socket MTH and TBB work stealing across node more efficient Qthread performance on multiple sockets is hampered by runtime overheads. WTI also increases as a result

## Relative speed to simulation

**Lessons learned:** 

- Task Size Matters: CPUs prefer fine-grained tasks GPUs prefer coarse-grained tasks
- Communication / disjoint memories
- Task runtime prediciton in irregular applications is a complex problem